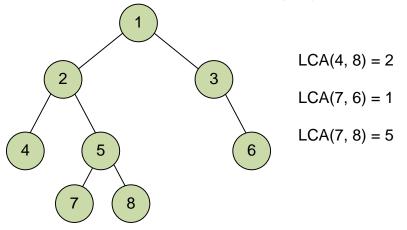
LCA - Least Common Ancestor

Let T be a rooted tree with *n* vertices. For two vertices *u* and *v*, you must find their closest common ancestor (the Least Common Ancestor). The procedure for finding such an ancestor will be denoted by LCA(u, v).

For example, if *u* is the ancestor of *v*, then LCA(u, v) = u.



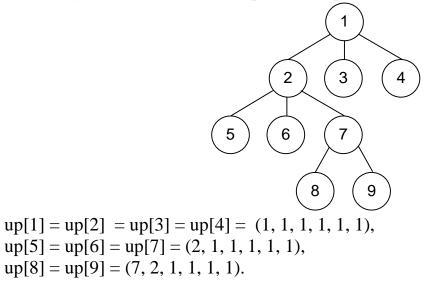
LCA. Binary lifting method

Let's find for each vertex its 1st, 2nd, 4th, 8th, ... ancestor. Store the results into the array *up*, where up[*i*][*j*] is equal to the 2^{j} - th ancestor of the vertex *i* ($1 \le i \le n$, $0 \le j \le \lceil \log_2 n \rceil$). If the 2^{j} -th ancestor of the vertex *i* does not exist, then set up[*i*][*j*] equal to the root of the tree.

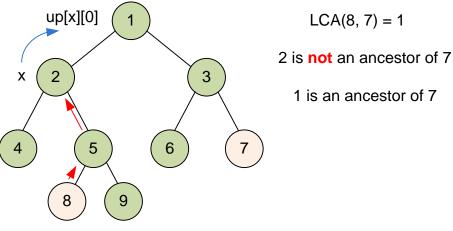
For each vertex v of the tree, compute the input d[v] and the output f[v] timestamps. They will be needed to determine in O(1) whether one vertex is an ancestor of another.

The preprocessing described is performed in $O(n\log_2 n)$.

Example. Consider the tree with n = 9 vertices. Let $l = \lceil \log_2 9 \rceil = 4$. Then each up[*i*] is an array of 5 elements (from up[*i*][0] to up[*i*][4]).



The query is to find the smallest common ancestor of vertices *a* and *b*. First, let's check if *a* is not an ancestor of *b*. Also, check if *b* is an ancestor of *a*. Otherwise, we'll lift the ancestors of the vertex *a* until we find the highest (closest to the root) vertex that is not yet an ancestor of *b* (not necessarily direct). That is, it will be a vertex *x* such that *x* is not an ancestor of *b*, but up[*x*][0] is already an ancestor of *b*. The query is executed in $O(\log_2 n)$ time.



E-OLYMP 2317. LCA offline (Easy) Find element in the tree.

Store the tree in the adjacency list g. Declare timestamps arrays d and f when traversing the tree with dfs. Declare an auxiliary array of ancestors up.

```
vector<vector<int> > g;
vector<int> d, f;
vector<vector<int> > up;
vector<pair<int, int> > Query;
char op[20];
```

Start the depth first search from the vertex v. The ancestor of v is the vertex p. Let the root of the tree be the vertex with number 1.

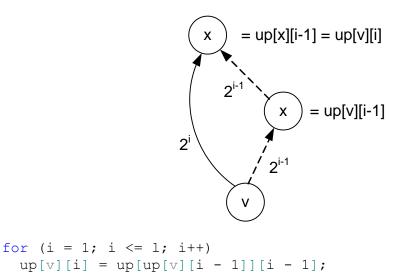
```
void dfs(int v, int p = 1)
{
    int i, to;
    d[v] = time++;
```

The immediate ancestor of *v* is *p*.

up[v][0] = p;

To find the 2^{i} -th ancestor of the vertex v, first find the 2^{i-1} -th ancestor of the vertex v, that equals to x = up[v][i - 1]. Then find the 2^{i-1} -th ancestor of the vertex x, that equals to

```
up[v][i] = up[x][i-1] = up[up[v][i-1]][i-1]
```



```
Continue dfs. Iterate over the vertices to, that can be reached from v.
```

```
for (i = 0; i < g[v].size(); i++)
{
   to = g[v][i];</pre>
```

If to is not an ancestor of v, then continue the search from the vertex to.

```
if (to != p) dfs(to, v);
}
f[v] = time++;
}
```

The function *Parent* returns 1 if *a* is the ancestor of *b*.

```
int Parent(int a, int b)
{
    return (d[a] <= d[b]) && (f[a] >= f[b]);
}
```

Function *LCA* returns the least common ancestor of the vertices *a* and *b*.

```
int LCA(int a, int b)
{
    if (Parent(a, b)) return a;
    if (Parent(b, a)) return b;
    for (int i = 1; i >= 0; i--)
        if (!Parent(up[a][i], b)) a = up[a][i];
    return up[a][0];
}
```

The main part of the program. Read the input data.

```
scanf("%d", &n);
g.resize(n + 1);
for (i = 0; i < n; i++)
{
   scanf("%s %d %d\n", op, &a, &b);
```

For the case of ADD query, add an edge to the tree. When a GET query is made, save its parameters in the *Query* array.

```
if (op[0] == 'A') { g[a].push_back(b); g[b].push_back(a); }
else Query.push_back(make_pair(a, b));
}
d.resize(n + 1); f.resize(n + 1); up.resize(n + 1);
```

Compute $l = \lceil \log_2 n \rceil$. Initialize an array *up*.

```
l = 1;
while ((1 << l) <= n + 1) l++;
for (i = 0; i <= n; i++) up[i].resize(l + 1);</pre>
```

Run the depth first search from the vertex 1.

dfs(1);

Compute and print the answers for the queries of type GET.

```
for (i = 0; i < Query.size(); i++)
printf("%d\n", LCA(Query[i].first, Query[i].second));</pre>
```